## **PROBLEM SET 9**

1. Let G be a graph whose edges are assigned lengths. Let t be the length of the shortest spanning tree of G. Let s be the length of the shortest path which visits every vertex.

**1.a** Show that  $t \leq s$ .

**1.b** Show that  $s \leq 2t$ .

Computing s is known as the traveling salesman problem, and is famously difficult. Nonetheless, this argument shows that it is easy to get a crude approximation to s.

**2.** Give an example of a graph G, with lengths assigned to edges with the following property: For any spanning tree T of G, there is some pair of vertices so that the distance from u to v in T is  $\geq 100$  times larger than the distance between u and v in G.

**3.** We showed in class that, if a graph has all edge lengths distinct, then there is a unique minimal spanning tree. Give an example of a graph where all edges are of distinct lengths, but the second and third best spanning tree have the same length.

4. The point of this problem is to prove a lemma we used in the course of verifying the Fürer-Raghavachari algorithm. Let G be a graph, T as spanning tree of G, and k the maximum degree of any vertex in T. Let S be the set of degree k vertices of T and B a subset of the degree k - 1 vertices of T; write s for |S| and b for |B|. Suppose that there are no edges of G joining distinct components of  $T \setminus (S \cup B)$ .

The result we will be establishing is that any spanning tree T' of G has a vertex of degree  $\geq k-1$ .

**4.a** Show that there are at most s + b - 1 edges of T connecting one vertex in  $S \cup B$  to another vertex in  $S \cup B$ .

**4.b** Show that there are at least (k-2)s + (k-3)b + 2 edges connecting a vertex in  $S \cup B$  to a vertex not in  $S \cup B$ .

Let H be the graph formed from G by deleting any edge of G which has one or more endpoints in  $S \cup B$ . So every vertex in  $S \cup B$  forms a connected component in H.

**4.c** Show that *H* has at least (k-1)s + (k-2)b + 2 connected components, including the isolated vertices from  $S \cup B$ .

**4.d** Now, let T' be any spanning tree of G. Show that there must be at least (k-1)s+(k-2)b+1 edges of T' with one more endpoints in  $S \cup B$ .

**4.e** Show that there is a vertex in  $S \cup B$  with T'-degree  $\geq k - 1$ .