



An Overview of Maximum Weight Graph Matching



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Maximum Weight Matching

For an undirected weighted graph, $G = (V, E, w)$:

• **Matching** : set of **endpoint disjoint** edges
 $M \subseteq E$

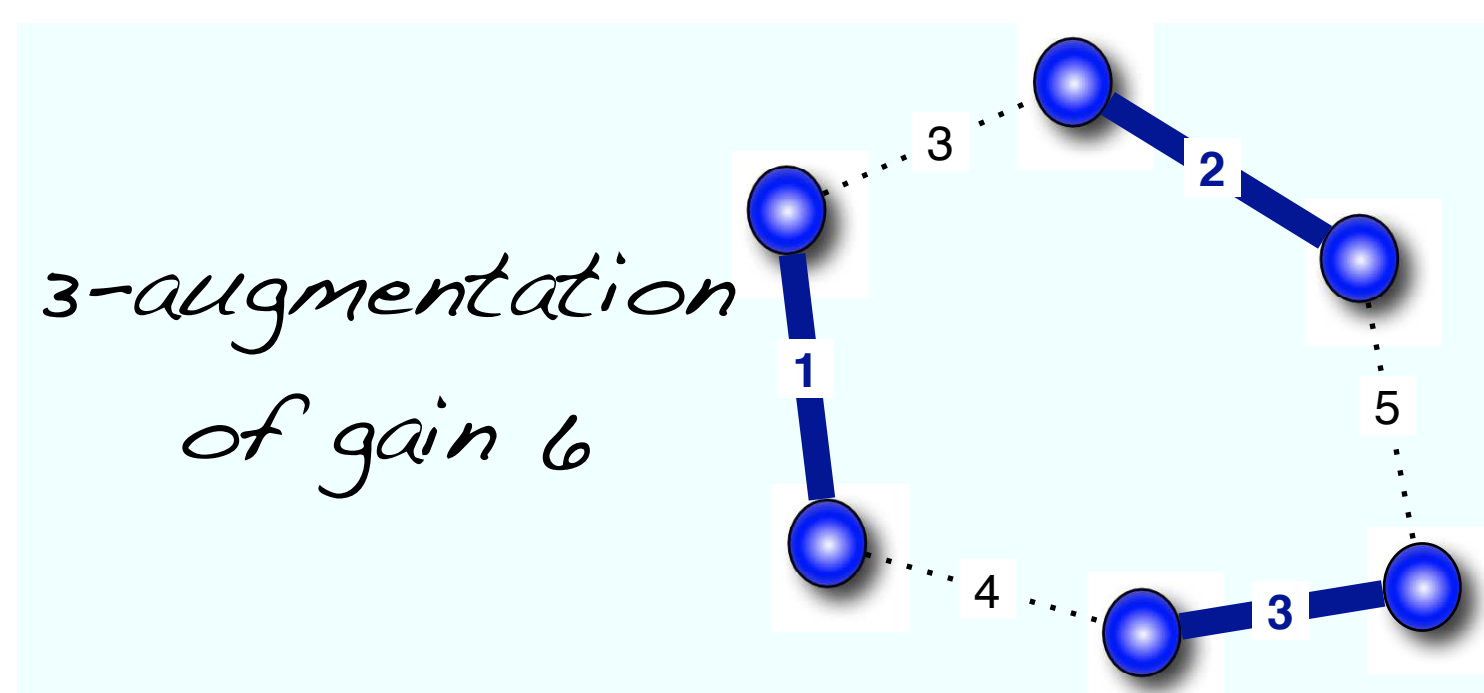
• **Weight of a Matching** :

$$w(M) = \sum_{e \in M} w(e)$$

Want to find **maximum weight matching**, M^* .

Augmentations and Gain

• **Augmentation** : Path or cycle in G with alternating matched and unmatched edges



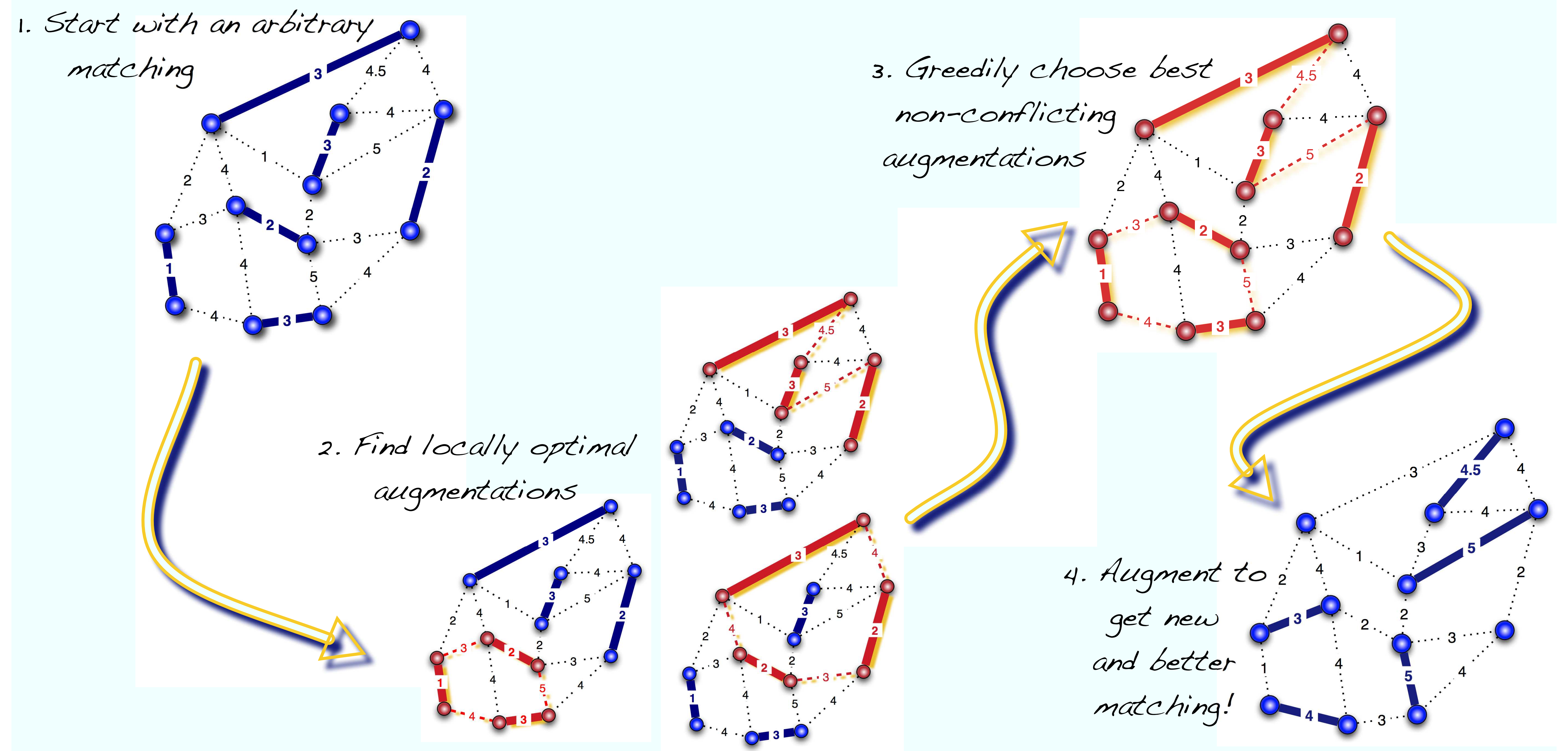
• **Gain of Augmentation** : $g(aug) = w(unmatched) - w(matched)$

• **k-augmentation** : Augmentation containing at most k unmatched edges

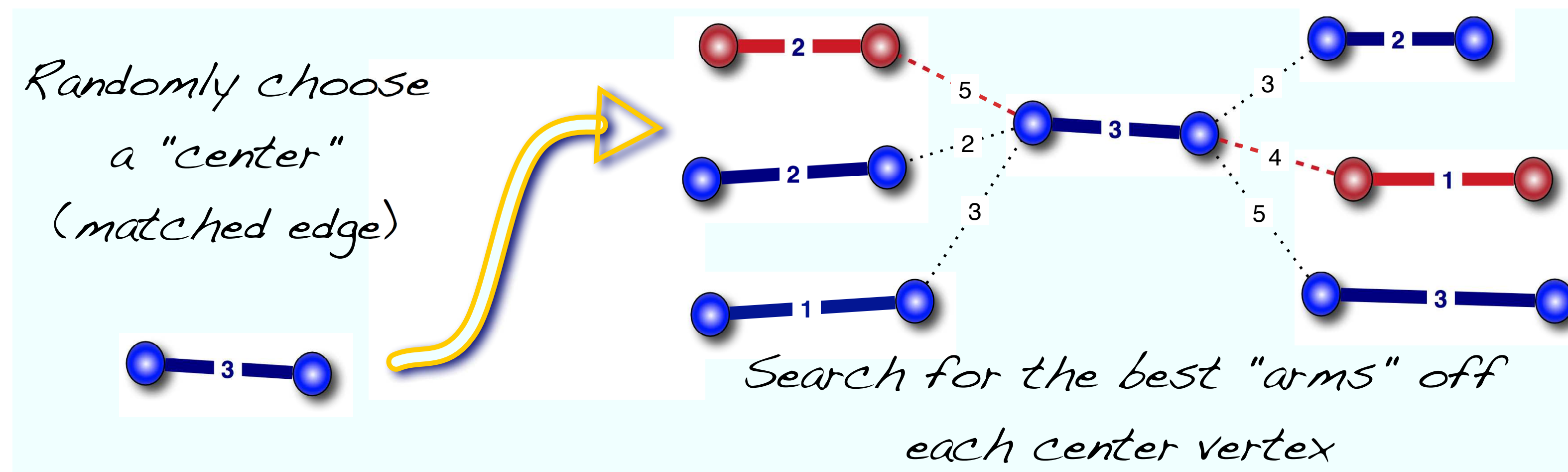
Approximation Theorem

If a matching has **no positive gain k-augmentation**, it is within $\frac{k}{k+1}$ of optimal.

Overview of Approximation Algorithm



Finding Locally Optimal Augmentations



Issues to Consider

- Sorting and updating lists of long arms
- Avoiding self-conflicting augmentations
- Finding cycles (or finding higher gain paths)