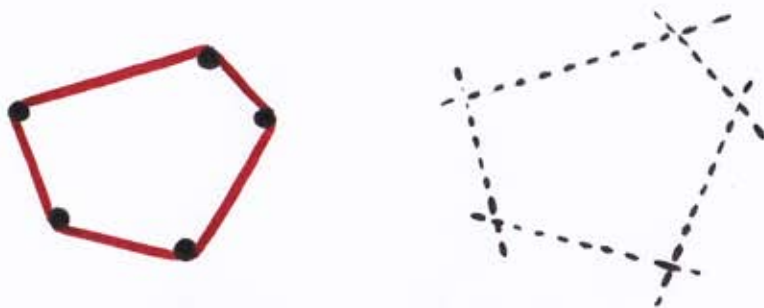


Polytopes

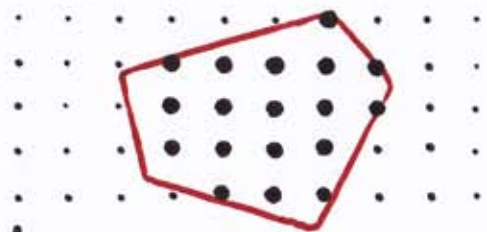
Generalization of polygon (2 dim.), polyhedron (3 dim.), etc. to \mathbb{R}^n .

Definition 1: Convex hull of finitely many points in \mathbb{R}^n .

Definition 2: Bounded region formed by intersection of finitely many half-spaces.



We want to count integer points:



Objects encoded by integer points: network flows, contingency tables, magic squares, Latin squares, knapsack packings, matchings...

Contingency tables

A contingency table is a nonnegative integer matrix.

EXAMPLE

	Black	Brown	Red	Blond	Σ	
Brown	68	119	26	7	220	row margins
Blue	20	84	17	94	215	
Hazel	15	54	14	10	93	
Green	5	29	14	16	64	
Σ	108	286	71	127	592	column margins

Cross-tab of eye, hair color.

$\Pi(R, C)$: polytope of nonneg. matrices with given margins
 $R = (r_1, \dots, r_m)$, $C = (c_1, \dots, c_n)$. This is the transportation polytope.

Applications:

- Cross-tabulation of categorical variables in a sample; significance testing.
- Program for moving goods from sites of supply to sites of demand; flow optimization.

Special- and generalizations:

- Magic squares
- Flows in arbitrary graphs
- Capacitated flows
- Higher-order tensors
- Latin squares / rook arrays

Counting integer points

Pick's Theorem (1899).

If P is a polygon with vertices in \mathbb{Z}^2 , then

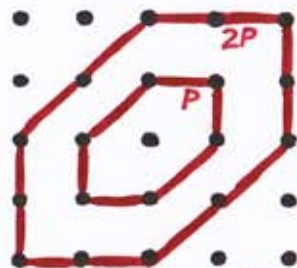
$$\text{Area}(P) = I + \frac{1}{2}B - 1.$$

integer points
in interior

integer points
on boundary

Ehrhart's Theorem (1967).

If P is a polytope with vertices in \mathbb{Z}^n , then there is a polynomial $l_P(t)$ which gives $|tP \cap \mathbb{Z}^n|$ whenever t is a nonnegative integer.



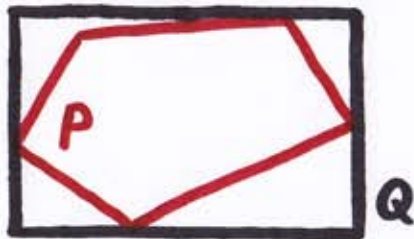
Ex.: $l_P(t) = 3t^2 + 3t + 1$

Implementation:

- Polynomial time when n fixed (Barvinok 1994; LattE)
- #P-complete in general

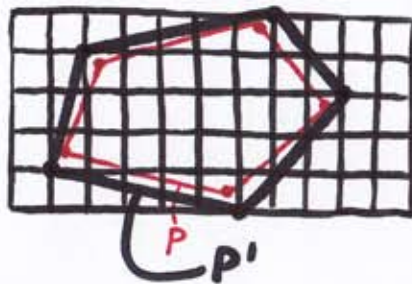
Coping in general dimension

Monte Carlo ("dart-throwing"):



- Choose points from $Q \cap \mathbb{Z}^n$ uniformly
- Observe $\frac{|P \cap \mathbb{Z}^n|}{|Q \cap \mathbb{Z}^n|}$

MC + dynamic programming:



- Replace P by a fixed-resolution simulacrum P' whose integer points can be counted iteratively in polynomial time
- Use iterates to throw darts at P'

MC + Markov chains:

- Sample $P \cap \mathbb{Z}^n$ by a random walk
- Must show good mixing
- Near-uniform sampling \leftrightarrow appx. counting
(Jerrum, Valiant, Vazirani)

Entropy methods ...

Independence models

Boxes are nice: If X is a uniform random integer point, then for $1 \leq j_1 < j_2 < \dots < j_r \leq n$ and $a_1, a_2, \dots, a_r \in \mathbb{Z}$,

$$\Pr \left[\bigwedge_{i=1}^r X_{j_i} = a_i \right] = \prod_{i=1}^r \Pr[X_{j_i} = a_i]. \quad (\heartsuit)$$

Might (\heartsuit) hold approximately for other polytopes, when $r \ll \dim P$?

(Inspiration: Phenomena in high-dimensional convex geometry...)

Definition.

An independence model is a random vector $X = (X_1, \dots, X_n)$ supported on \mathbb{Z}^n which satisfies (\heartsuit) . In other words, the coordinates are independent.

An independence model for a simplex truncated by a cube

Let $P = \{(x_1, \dots, x_n) \in [0, 1]^n : x_1 + \dots + x_n = r\}$.

If Y is uniformly distributed on $P \cap \mathbb{Z}^n$, then Y_j is 0-1 Bernoulli and has mean r/n .

Let X be the independence model with $X_j \sim Y_j$, $1 \leq j \leq n$.

Then X hits every point of $P \cap \mathbb{Z}^n$ at rate $e^{-nh(r/n)}$,
where $h(p) = -p \ln p - (1-p) \ln(1-p)$.

Also, $\Pr(X \in P) \approx \sqrt{\frac{n}{2\pi r(n-r)}}$, so

$$|P \cap \mathbb{Z}^n| \approx e^{nh(r/n)} \sqrt{\frac{n}{2\pi r(n-r)}}.$$

We recover the entropy formula for $\binom{n}{r}$:

Proposition. $\ln \binom{sn}{sr} = sn \cdot h\left(\frac{r}{n}\right) - \Theta(\ln s)$.

Entropy

Entropy is a statistic associated to a random variable:

$$H[X] = \sum_{x \in \text{supp } X} \text{Pr}[X=x] \ln \frac{1}{\text{Pr}[X=x]}.$$

(Assume X discrete, but $H[X]$ can still be ∞ .)

Basic properties:

- $H[X]$ is strictly concave w.r.t. the p.m.f. of X .
(Thus $H[X] \leq \ln |\text{supp } X|$, achieved by the uniform distr.)
- $H[X, Y] \leq H[X] + H[Y]$,
with equality iff. X, Y independent.

Maximum entropy principle:

Given partial information about a random variable X , Jaynes proposed the maximum-entropy distribution among all candidates as the likeliest...

The maximum-entropy independence model

Let $P = \{x \in \mathbb{R}^n : x \geq 0, Ax = b, (x \leq k)\}$.

$A \in \mathbb{R}^{m \times n}$
of rank m

$b \in \mathbb{R}^m$

$k \in (\mathbb{R} \cup \{\infty\})^n$

Let $X = (X_1, \dots, X_n)$ be a random vector with max. entropy subject to:

- $\text{supp } X \subseteq \mathbb{Z}_{\geq 0}^n$
- $\mathbb{E}[AX] = b$ (equiv.: $\mathbb{E}[X] \in P$)

Properties:

- X exists and has unique distribution
- X is an independence model (called the MEIM)
- X has constant mass $e^{-H[X]}$ on points of $P \cap \mathbb{Z}^n$
- Coordinates of X are (truncated) geometric:

$$\Pr[X_j = t] = p_j q_j^t \quad (t=0, 1, 2, \dots, k_j)$$

for some parameters $p_j \in [0, 1]$, $q_j \in [0, \infty]$

(Note: $q_j = \infty$ requires special interpretation; also, $q_j \in [0, 1)$ if $k_j = \infty$)

Truncated geometric distribution

For $\kappa \in \mathbb{Z}_{>0} \cup \{\infty\}$, $^{0 \leq x \leq \kappa}$, we say $X \sim TG(x; \kappa)$ if

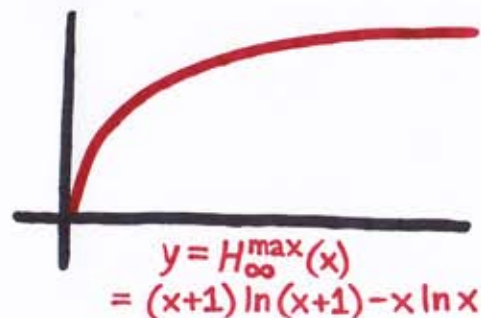
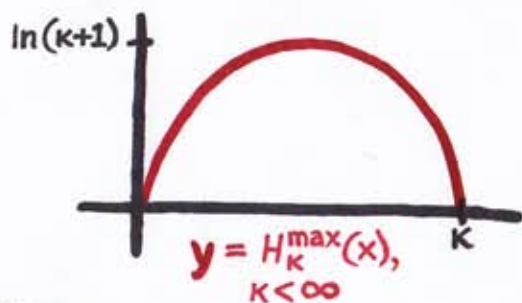
$$\Pr(X=t) = pq^t \quad \text{for } t=0,1,\dots,(\kappa),$$

where $p = p(x; \kappa)$, $q = q(x; \kappa)$ satisfy

$$\begin{cases} 1 = p(1+q+q^2+\dots+q^\kappa) \\ x = p(q+2q^2+\dots+\kappa q^\kappa) \end{cases} .$$

This is the max. ent. distr. on $\{0,1,\dots,\kappa\}$ with mean x .

Let $H_\kappa^{\max}(x)$ denote its entropy.



Properties:

- Nice formulas for p, q, H_κ^{\max} when $\kappa=1$ or $\kappa=\infty$
- $H_\kappa^{\max}(x)$ is strictly concave in x
- $H_\kappa^{\max}(x) = -[\ln p + x \ln q]$
- $(H_\kappa^{\max})'(x) = -\ln q$

Using the MEIM to count integer points

Let X be the MEIM for P .

$$\text{Then } |P \cap \mathbb{Z}^n| = \underbrace{e^{H[X]}}_{\text{This is the max. of a strictly concave function on a polytope, so it can be computed in polynomial time.}} \underbrace{\text{Pr}[X \in P]}_{\text{We want to estimate this! Equivalently, to estimate } \text{Pr}[AX=b].}$$

The "I-bound":

Theorem. Let $P = \{x \in \mathbb{R}^n: x \geq 0, Ax = b\}$, where $A = \begin{pmatrix} \bar{a}_1 & \bar{a}_2 & \dots & \bar{a}_n \end{pmatrix} \in \mathbb{R}^{m \times n}$ is of rank m and $b \in \mathbb{R}^m$.

Let X be the MEIM associated to P , and $q_j = q(\mathbb{E}[X_j]; \kappa)$, $1 \leq j \leq n$.

$$\text{Then } |P \cap \mathbb{Z}^n| \leq e^{H[X]} \min_{\substack{a_{j_1}, \dots, a_{j_m} \\ \text{lin. indep.}}} \prod_{i=1}^m \underbrace{(1 - q_{j_i})}_{\text{or } \frac{1}{1 + \mathbb{E}(X_{j_i})}}.$$

The Littlewood-Offord problem

L., O. asked: How many subsets of $\{a_1, \dots, a_n\}$ can have equal sum?

When $a_1, \dots, a_n \in \mathbb{R} \setminus \{0\}$, the answer is $\binom{n}{\lfloor n/2 \rfloor}$ (Erdős, 1945).

The case of $a_1, \dots, a_n \in \mathbb{R}^m \setminus \{0\}$ was addressed by Halász:

Theorem (Halász, 1977).

Let $a_1, \dots, a_n \in \mathbb{R}^m$. Let $\varepsilon_1, \dots, \varepsilon_m$ be independent symmetric Bernoulli random variables, and let $S := \varepsilon_1 a_1 + \dots + \varepsilon_n a_n$.

Define $\text{conc}_1(S) := \max_{y \in \mathbb{R}^m} \Pr(|S - y| < 1)$.

Suppose there exists $\delta > 0$, such that for any $|e|=1$ one can select at least δn vectors a_k with $|\langle a_k, e \rangle| \geq 1$. Then

$$\text{conc}_1(S) \leq c(\delta, m) n^{-m/2},$$

where $c(\delta, m)$ depends only on δ and m .

By a scaling limit argument, we can replace each 1 by an arbitrarily small $\varepsilon > 0$ and pass to a "point concentration" version.

The "E-bound"

Theorem. Let $P = \{x \in \mathbb{R}^n: x \geq 0, Ax = b\}$, where

$A = \begin{bmatrix} \bar{a}_1 & \bar{a}_2 & \dots & \bar{a}_n \end{bmatrix} \in \mathbb{R}^{m \times n}$ and $b \in \mathbb{R}^m$. Let X be the MEIM,

and suppose $2 \leq \mathbb{E}[X_j] < N$ for all $j=1, \dots, n$.

Suppose $n = pm$ and $\hat{a}_i = \hat{a}_{m+i} = \hat{a}_{2m+i} = \dots = \hat{a}_{(p-1)m+i}$ for $i=1, \dots, m$, where $\{a_1, \dots, a_m\}$ is a basis for \mathbb{R}^m .

Then for fixed N, m , we have

$$|P \cap \mathbb{Z}^n| \leq (1+o(1)) e^{H(X)} \prod_{i=1}^m \left[\frac{\pi}{6} \sum_{t=1}^p (|\mathbb{E}[X_{(t-1)m+i}] + 1|^2 - 1) \right]^{-1/2}$$

as $p \rightarrow \infty$.

Key lemmas:

- If X_1, \dots, X_p are indep. discrete RVs s.th. $\text{conc}(X_j) \leq \frac{1}{N_j} \forall j$, then $\text{conc}(X_1 + \dots + X_p)$ is maximized by X_j uniformly distributed on $\{0, 1, \dots, N_j - 1\}$.
- Bender's local limit theorem for RVs with log-concave p.m.f.

$$\text{conc}(Y) := \max_{y \in \text{supp } Y} \Pr[Y=y]$$

The "H-bound"

Theorem (executive summary).

For every integer $m \geq 1$ and real $\varepsilon > 0$, there exists $\delta = \delta(m, \varepsilon) > 0$, such that the following holds:

Let P be a polytope in the form $P = \{x \in \mathbb{R}^n : x \geq 0, Ax = b\}$, where $A \in \mathbb{Z}^{m \times n}$, $b \in \mathbb{R}^m$, and the columns of A span \mathbb{R}^m p times ($p \geq 1$). Let X be the MEIM for P . Then if $E[X_j] \geq \varepsilon$ for $1 \leq j \leq n$, we have

$$|P \cap \mathbb{Z}^n| \leq \delta p^{-m/2} e^{H[X]}.$$

The full result includes explicit constants:

$$|P \cap \mathbb{Z}^n| \leq (p^{-m/2} \cdot C + (C')^p) e^{H[X]},$$

where C, C' are given in terms of $\min_{1 \leq t \leq p} E[X_{(t-1)m+i}]$, $i = 1, \dots, m$ (where the first p blocks of m columns of A are assumed to be bases).

Numerical examples

- Contingency tables with $R = (220, 215, 93, 64)$,
 $C = (108, 286, 71, 127)$:

$$\begin{array}{ll} H\text{-bound: } 8.01 \times 10^{26} & I\text{-bound: } 7.14 \times 10^{18} \\ \text{Actual count: } 1.23 \times 10^{15} & \end{array}$$

- $3 \times 3 \times 3$ tables with layer sums equal to 20:

$$\begin{array}{ll} H\text{-bound: } 3.66 \times 10^{20} & I\text{-bound: } 7.00 \times 10^{19} \\ \text{Actual count: } 6.43 \times 10^{14} & \end{array}$$

- Integer points of a simplex $2x_1 + 11x_2 + 18x_3 + \dots + 14x_{25} = 5000$:

$$\begin{array}{ll} H\text{-bound: } 2.00 \times 10^{44} & I\text{-bound: } 1.07 \times 10^{44} \\ E\text{-bound: } 1.04 \times 10^{44} & \text{Actual count: } 8.57 \times 10^{42} \end{array}$$

- Integer points of the simplex $x_1 + x_2 + \dots + x_{10000} = 100$:

$$\begin{array}{l} H\text{-bound: } 1.774 \times 10^{242} \\ \text{Actual count: } 1.755 \times 10^{242} \end{array}$$

Bounded contingency tables

The truncated transportation polytope:

$$\Pi_K(R, C) := \left\{ X \in \mathbb{R}_{\geq 0}^{m \times n} : \sum_{j=1}^n x_{ij} = r_i \quad (i=1, \dots, m), \right. \\ \left. \sum_{i=1}^m x_{ij} = c_j \quad (j=1, \dots, n), \right. \\ \left. x_{ij} \leq k_{ij} \text{ for all } i, j \right\},$$

where

$$R = (r_1, \dots, r_m) \in \mathbb{Z}_{\geq 0}^m, \quad C = (c_1, \dots, c_n) \in \mathbb{Z}_{\geq 0}^n,$$

$$N = r_1 + \dots + r_m = c_1 + \dots + c_n, \quad \text{and } K = (k_{ij}) \in (\mathbb{Z}_{\geq 0} \cup \{\infty\})^{m \times n}.$$

The integer points of $\Pi_K(R, C)$ are K -bounded contingency tables.

Let $T_K(R, C)$ be the # of such tables.

Let $I_K(R, C)$ be the independence estimate

$$\left(\begin{array}{l} \# \text{ tables w/ sum} \\ \text{of entries} = N \end{array} \right) \cdot \left(\begin{array}{l} \text{fraction of those} \\ \text{tables having row} \\ \text{sums } r_1, \dots, r_m \end{array} \right) \cdot \left(\begin{array}{l} \text{fraction having} \\ \text{column sums} \\ c_1, \dots, c_n \end{array} \right).$$

Counting contingency tables via permanents

Theorem (Barvinok).

Let $R \in \mathbb{Z}_{\geq 0}^m$, $C \in \mathbb{Z}_{\geq 0}^n$, $|R| = |C| = N$ as before.

Let $W = (w_{ij}) \in \mathbb{R}^{m \times n}$. Let $\Gamma = (\gamma_{ij})$ be an $m \times n$ matrix whose entries are independent exponential random variables of mean 1. Let $A = A(\Gamma)$ be the $N \times N$ matrix obtained by replacing each entry (i, j) of Γ by an $r_i \times c_j$ block with all entries equal to $w_{ij} \gamma_{ij}$.

Let each contingency table X with margins (R, C) be counted with weight $w(X) = \prod_i \prod_j w_{ij}^{x_{ij}}$.

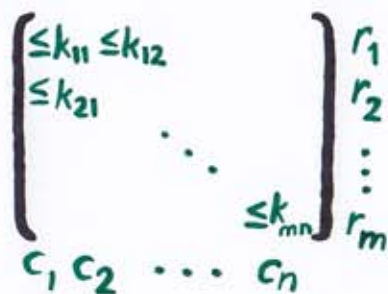
Then the total weight of all such tables is

$$T(R, C; W) := \frac{\mathbb{E}[\text{per } A]}{r_1! \cdots r_m! c_1! \cdots c_n!}.$$

Using matrix scaling and bounds on the permanent of a doubly stochastic matrix, we can compute $\mathbb{E}[\text{per } A]$ to within a factor $N^{O(m+n)}$ in polynomial time.

Adapting Barvinok's theorem to bounded tables

Assume
 $k_{ij} < \infty$
 $\forall i, j.$



$(m \times n)$ bounded table

bijection



$((m+n) \times (mn))$ table with entries \bullet counted at weight 1, others at weight 0

$$T_K(R, C) = T(\mathcal{R}, \mathcal{C}; W)$$

where $\mathcal{R} = (r_1, \dots, r_m, \tilde{c}_1, \dots, \tilde{c}_n)$, $\mathcal{C} = (k_{11}, \dots, k_{1n}, k_{21}, \dots, k_{2n}, \dots, k_{m1}, \dots, k_{mn})$,

$$\tilde{c}_j = \sum_{i=1}^m k_{ij} - c_j.$$

$T_K(R, C)$ is approximately log-concave w.r.t. R, C ; that is,

$\ln T_K(R, C)$ can be written as the sum of a concave function $f(R, C)$ and a function bounded by $O(m+n + \sum_i \ln r_i + \sum_j \ln c_j)$.

Exact and approximate generating functions

From the last slide: $\ln T_K(R, C) = \underbrace{f(R, C)}_{\text{concave approximation from above}} + \underbrace{\text{remainder}}_{O(m+n+\sum \ln r_i + \sum \ln c_j)}$.

Let

$$G(\vec{x}, \vec{y}) = G_K(\vec{x}, \vec{y}) := \prod_{i=1}^m \prod_{j=1}^n [1 + x_i y_j + (x_i y_j)^2 + \dots + (x_i y_j)^{k_{ij}}]$$

$$= \sum_R \sum_C T_K(R, C) x^R y^C.$$

Let $\tilde{G}(\vec{x}, \vec{y}) = \sum_R \sum_C e^{f(R, C)} x^R y^C.$

Upper bound: $T_K(R, C) \leq \inf_{x_i, y_j > 0} \frac{G(\vec{x}, \vec{y})}{x^R y^C}.$

Lower bound:

Given R, C , may choose x_*, y_* so that $e^{f(R, C)} x_*^R y_*^C$ is the largest term in the expansion of $\tilde{G}(x_*, y_*)$.

Thus $\frac{\tilde{G}(x_*, y_*)}{x_*^R y_*^C} \leq e^{f(R, C)} \cdot (\# \text{ nonzero terms of } \tilde{G}),$

so $\inf_{x_i, y_j > 0} \frac{G(\vec{x}, \vec{y})}{x^R y^C} \leq T_K(R, C) \cdot \underbrace{(\# \text{ nonzero terms of } \tilde{G})}_{\leq \prod_i (1 + \sum_j k_{ij}) \cdot \prod_j (1 + \sum_i k_{ij})} \cdot \underbrace{e^{\text{remainder}}}_{e^{O(m+n+\sum \ln r_i + \sum \ln c_j)}}.$

Asymptotic formulas for $\ln T_K(R, C)$, part 1

Theorem. Let d be fixed, and let k_{ij} vary between fixed positive bounds. Then

$$\ln T_K(R, C) = \ln \left(\inf_{x_i, y_j > 0} \frac{G(x, y)}{x^R y^C} \right) + O(\max\{m, n\} \ln \max\{m, n\}) \quad (\diamond)$$

uniformly for R, C satisfying

$$\max\{r_1, \dots, r_m\} \leq \frac{d}{m} (r_1 + \dots + r_m),$$

$$\max\{c_1, \dots, c_n\} \leq \frac{d}{n} (c_1 + \dots + c_n).$$

Moreover, the first term in (\diamond) is $e^{\Omega(mn)}$, so

$$\ln T_K(R, C) \sim \ln \inf_{x_i, y_j > 0} \frac{G(x, y)}{x^R y^C}$$

as $m, n \rightarrow \infty$.

Asymptotic formulas for $\ln T_K(R, C)$, part 2

The MEIM for $\Pi_K(R, C)$ is a matrix $X = (X_{ij})$ of indep. TG random variables with mean $Z \in \Pi_K(R, C)$ maximizing $\sum_i \sum_j H_{k_{ij}}^{\max}(z_{ij})$. ($= H[X]$)

We know that X is uniform on the int. pts. of $\Pi_K(R, C)$, so

$$H[X] \geq \ln T_K(R, C).$$

This upper bound on $\ln T_K(R, C)$ turns out to be asymptotic to the true value, thanks to

Lemma.
$$\ln \left(\inf_{x_i, y_j > 0} \frac{G(x, y)}{x^R y^C} \right) = \max_{Z \in \Pi_K(R, C)} \sum_i \sum_j H_{k_{ij}}^{\max}(z_{ij}).$$

Hence

Theorem.
$$\ln T_K(R, C) \sim \max_{Z \in \Pi_K(R, C)} \sum_i \sum_j H_{k_{ij}}^{\max}(z_{ij})$$

in the asymptotic régime of the last theorem.

Detecting correlation of margins, part 1

We say that margins R, C are positively (or negatively) correlated if $T_K(R, C) > I_K(R, C)$ (or $T_K(R, C) < I_K(R, C)$).

We study the case of K constant: $k_{ij} = \kappa$ for all i, j , where $\kappa \in \mathbb{Z}_{>0} \cup \{\infty\}$. Suppose neither R nor C is a constant vector.

Barvinok observed that, as $s \rightarrow \infty$, the margins

$$R^{(s)} = (sr_1, \dots, sr_m, \underbrace{sr_1, \dots, sr_m}_{s \text{ repetitions}}, \dots, sr_1, \dots, sr_m),$$

$$C^{(s)} = (sc_1, \dots, sc_n, \underbrace{sc_1, \dots, sc_n}_{s \text{ repetitions}}, \dots, sc_1, \dots, sc_n)$$

are asymptotically correlated with positive sign and

$$\frac{T_K(R^{(s)}, C^{(s)})}{I_K(R^{(s)}, C^{(s)})} = e^{\Omega(s^2)} \quad \text{when } \kappa = \infty, \text{ but}$$

with negative sign and

$$\frac{I_K(R^{(s)}, C^{(s)})}{T_K(R^{(s)}, C^{(s)})} = e^{\Omega(s^2)} \quad \text{when } \kappa = 1.$$

What does the transition look like?

Detecting correlation of margins, part 2

There are positively correlated margins whenever $\kappa \geq 2$:

Theorem. Let $R \in \mathbf{Z}_{>0}^m$, $C \in \mathbf{Z}_{>0}^n$, and $\kappa \in \{2, 3, 4, \dots\}$.

Then there exists $\delta = \delta(\kappa) \in (0, 1)$, such that if (R, C) satisfy

$$\left(\max_i r_i\right) \left(\max_j c_j\right) < \delta \kappa N,$$

then

$$\lim_{s \rightarrow \infty} \frac{1}{s^2} \ln T_\kappa(R^{(s)}, C^{(s)}) \geq \lim_{s \rightarrow \infty} \frac{1}{s^2} \ln I_\kappa(R^{(s)}, C^{(s)}),$$

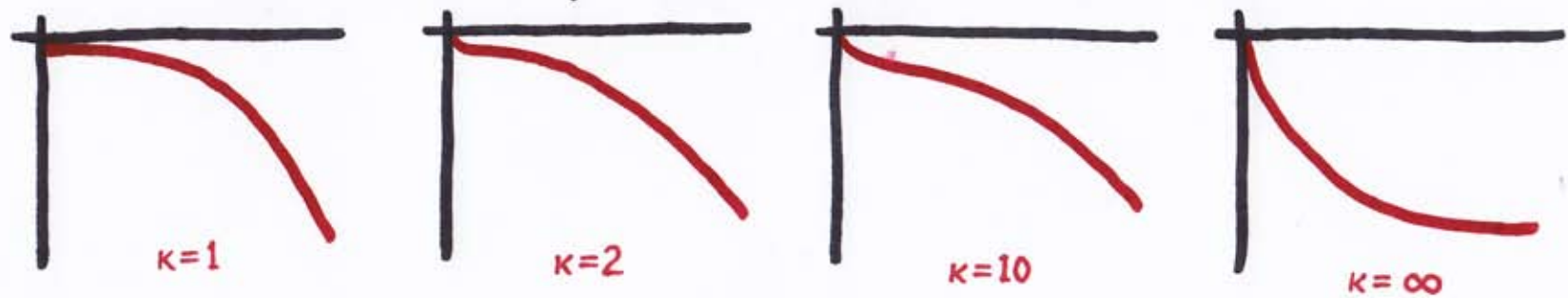
with strict inequality if neither R nor C is a constant vector.

Proof idea:

- Both $T_\kappa(R, C)$, $I_\kappa(R, C)$ can be estimated via entropy formulas.
- Compare four independence models with TG entries satisfying $E[X_{ij}] = \frac{N}{mn}$, $\frac{r_i}{n}$, $\frac{c_j}{m}$, $\frac{r_i c_j}{N}$. The entropies of the first two differ by more than the entropies of the last two. Thus (informally) learning R causes less surprise when we know C .
- Actual proof relies heavily on convexity of $\phi(x) := x^2 (H_\kappa^{\max})''(x)$.

Detecting correlation of margins, part 3

The decisive function $\phi(x) = x^2(H_\kappa^{\max})''(x)$:



Conjecture: Let $R \in \mathbb{Z}_{>0}^m$, $C \in \mathbb{Z}_{>0}^n$, and $\kappa \in \{2, 3, 4, \dots\}$.

Let δ_{cr} be the largest value of δ possible in the previous theorem. Then if (R, C) satisfy

$$\delta_{cr} \kappa n < r_i < (1 - \delta_{cr}) \kappa n \quad \text{for } i = 1, \dots, m,$$

$$\delta_{cr} \kappa m < c_j < (1 - \delta_{cr}) \kappa m \quad \text{for } j = 1, \dots, n,$$

the conclusion of the previous theorem holds with reversed sign.

Numerical evidence suggests:

$$\kappa = 2 \quad 4 \quad 6 \quad 8 \quad 10$$

$$\delta_{cr} \approx .05 \quad .11 \quad .14 \quad .15 \quad .16$$