Parametrizations of *k*-Nonnegative Matrices

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Factorizations

Cluster Algebras

Background

In 1999, Fomin and Zelevinsky studied totally nonnegative matrices.

They explored two questions:

- 1. How can totally nonnegative matrices be parameterized?
- 2. How can we test a matrix for total positivity?

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- 1. How can totally nonnegative matrices be parameterized?
- 2. How can we test a matrix for total positivity?

We will explore the same questions for k-nonnegative and k-positive matrices.

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Lemma

A matrix M is k-nonnegative if all column-solid minors of order k or less are nonnegative.

Factorizations

Chevalley generators

Loewner-Whitney Theorem: An invertible totally nonnegative matrix can be written as a product of e_i 's, f_i 's and h_i 's with nonnegative entries.

$$e_{i}(a) = \begin{bmatrix} 1 & 0 & \cdots & \cdots & 0 \\ \vdots & \ddots & 0 & \cdots & \cdots & \vdots \\ 0 & \cdots & 1 & a & \cdots & 0 \\ 0 & \cdots & 0 & 1 & \cdots & 0 \\ \vdots & \cdots & \cdots & \ddots & \ddots & \vdots \\ 0 & \cdots & \cdots & \cdots & 0 & 1 \end{bmatrix}, f_{i}(a) = \begin{bmatrix} 1 & 0 & \cdots & \cdots & 0 \\ \vdots & \ddots & 0 & \cdots & \cdots & 0 \\ 0 & \cdots & a & 1 & \cdots & 0 \\ \vdots & \cdots & \cdots & \ddots & \vdots \\ 0 & \cdots & \cdots & 0 & 1 \end{bmatrix}$$
$$h_{i}(a) = \begin{bmatrix} 1 & 0 & \cdots & \cdots & 0 \\ 0 & \ddots & 0 & \ddots & \vdots \\ \vdots & \ddots & a & \ddots & \vdots \\ \vdots & \ddots & a & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & 0 \\ 0 & \cdots & \cdots & 0 & 1 \end{bmatrix}$$

Row and Column Reductions

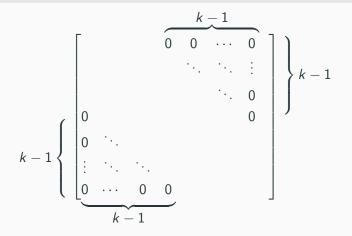
Lemma

If a matrix M is k-nonnegative, it can be reduced to have a k-1 "staircase" of 0s in its northeast and southwest corners while preserving k-nonnegativity.

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Generators

Theorem

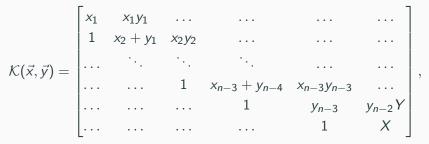
The semigroup of n - 1-nonnegative invertible matrices is generated by the Chevalley generators and the \mathcal{K} -generators.

Generators

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The semigroup of n - 1-nonnegative invertible matrices is generated by the Chevalley generators and the \mathcal{K} -generators.

The $\mathcal K\text{-}\mathsf{generators}$ have the following form.



$$Y = y_1 \cdots y_{n-3}$$

$$X = x_2 x_3 \cdots x_{n-3} + y_1 x_3 \cdots x_{n-3} + y_1 y_2 x_3 \cdots x_{n-3} + \dots + y_1 \cdots y_{n-4}.$$

$$\begin{split} \mathbf{e}_{j}(a) \cdot \mathcal{K}(\vec{x}, \vec{y}) &= \mathcal{K}(\vec{u}, \vec{v}) \cdot \mathbf{e}_{j+1}(b) \text{ where } 1 \leq j \leq n-2 \\ \mathbf{e}_{n-1}(a) \cdot \mathcal{K}(\vec{x}, \vec{y}) &= h_{n}(b) \cdot \mathcal{K}(\vec{u}, \vec{v}) \cdot \mathbf{f}_{n-1}(c) \\ h_{j+2}(c) \cdot \mathbf{f}_{j+1}(a) \cdot \mathcal{K}(\vec{x}, \vec{y}) &= \mathcal{K}(\vec{u}, \vec{v}) \cdot \mathbf{f}_{j}(b) \cdot h_{j}(c) \text{ where } 1 \leq j \leq n-2 \\ \mathbf{f}_{1}(a) \cdot \mathcal{K}(\vec{x}, \vec{y}) \cdot h_{1}(c) &= \mathcal{K}(\vec{u}, \vec{v}) \cdot \mathbf{e}_{1}(c) \\ h_{j+1}(a) \cdot \mathcal{K}(\vec{x}, \vec{y}) &= \mathcal{K}(\vec{u}, \vec{v}) \cdot h_{j}(a) \text{ where } 1 \leq j \leq n-2. \end{split}$$

Generators

Theorem

The semigroup of n - 2-nonnegative upper unitriangular matrices is generated by the e_i 's and the T-generators.

The $\ensuremath{\mathcal{T}}\xspace$ -generators have the following form.

	[1	x_1	x_1y_1			•••]
		1	$x_2 + y_1$	<i>x</i> ₂ <i>y</i> ₂			
			÷.,	·	$ \begin{array}{c} \cdots \\ \cdots \\ \vdots \\ x_{n-3} + y_{n-4} \\ 1 \\ \cdots \\ \cdots \\ \cdots \end{array} $		
$\mathcal{T}(\vec{x},\vec{y}) =$		• • •		1	$x_{n-3} + y_{n-4}$	$x_{n-3}y_{n-3}$	
					1	У <i>п</i> —3	$y_{n-2}Y$
						1	X
	[1

$$Y = y_1 \cdots y_{n-3}$$

$$X = x_2 x_3 \cdots x_{n-3} + y_1 x_3 \cdots x_{n-3} + y_1 y_2 x_3 \cdots x_{n-3} + \dots + y_1 \cdots y_{n-4}.$$

$$\begin{aligned} \mathbf{e}_{j}(\mathbf{a}) \cdot \mathcal{T}(\vec{x}, \vec{y}) &= \mathcal{T}(\vec{u}, \vec{v}) \cdot \mathbf{e}_{j+2}(b) \text{ where } 1 \leq j \leq n-3 \\ \mathbf{e}_{n-2}(\mathbf{a}) \cdot \mathcal{T}(\vec{x}, \vec{y}) &= \mathcal{T}(\vec{u}, \vec{v}) \cdot \mathbf{e}_{1}(b) \\ \mathbf{e}_{n-1}(\mathbf{a}) \cdot \mathcal{T}(\vec{x}, \vec{y}) &= \mathcal{T}(\vec{u}, \vec{v}) \cdot \mathbf{e}_{2}(b) \end{aligned}$$

Alphabet
$$\mathcal{A} = \{1, 2, \dots, n-1, \mathcal{T}\}.$$

Let α be the word $(n-2) \dots 1(n-1) \dots 1.$

Alphabet $\mathcal{A} = \{1, 2, \dots, n-1, \mathcal{T}\}$. Let α be the word $(n-2) \dots 1(n-1) \dots 1$. The reduced words are:

$$w \in \begin{cases} w'\mathcal{T} & w'\alpha \text{ is reduced,} \\ w'(n-1)\mathcal{T} & w'\alpha \text{ is reduced,} \\ w'(n-2)\mathcal{T} & w'\alpha \text{ is reduced,} \\ w'(n-1)(n-2)\mathcal{T} & w'\alpha \text{ is reduced,} \\ w' & w' < \beta \text{ or } w' \text{ is incomparable to } \beta. \end{cases}$$

where w' does not involve \mathcal{T} .

Define V(w) to be the set of matrices which correspond to the reduced word w. (Then $V(w) = \{e_{w_1}(a_1)e_{w_2}(a_2)\cdots e_{w_k}(a_k)\}$.)

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Theorem

For reduced words u and w, if $u \neq w$ then $V(u) \cap V(w) = \emptyset$.

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Theorem

The poset on $\{V(w)\}$ given by the Bruhat order on reduced words $\{w\}$ is graded.

Conjecture

The closure of a cell $\overline{V(w)}$ is the disjoint union of all cells in the interval between \emptyset and V(w).

Cluster Algebras

k-initial minors

Definition

A *k*-initial minor at location (i, j) of a matrix X is the maximal solid minor with (i, j) as the lower right corner which is contained in a $k \times k$ box.

The set of all k-initial minors gives a k-positivity test!

[11	12	13	14	15	16		[11	12	13	14	15	16
21							1					
31												
41	42	43	44	45	46	,	41	42	43	44	45	46
51	52	53	54	55	56		51	52	53	54	55	56
61	62	63	64	65	66		61	62	63	64	65	66

4-initial minors

With total positivity tests, can "exchange" some minors for others.

Example

$$M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$$

Both $\{a, b, c, \det M\}$ and $\{d, b, c, \det M\}$ give total positivity tests.

Note

$$ad = bc + \det M$$

i.e. have a *subtraction-free* expression relating exchanged minors.

Definition

A seed is a tuple of variables $\tilde{\mathbf{x}}$ along with some exchange relations of the form

$$x_i x'_i = p_i (\tilde{\mathbf{x}} \setminus x_i)$$

which allow variable x_i to be swapped for a new variable x'_i .

- frozen variables: not exchangeable
- cluster variables: are exchangeable
- extended cluster: entire tuple $\tilde{\mathbf{x}}$
- cluster: only the cluster variables

A seed (plus all seeds obtained by doing chains of exchanges) generates a *cluster algebra*. Our p_i are always *subtraction-free*.

Example

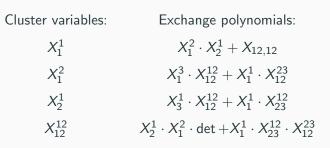
Initial seed: $\tilde{\mathbf{x}}$ is minors of *n*-initial minors test. *Corner minors* (lower right corner on bottom or right edge) are frozen variables. There is a rule for generating the exchange relations for all other variables.

Subtraction-freeness means that any seed reachable from the initial one gives a different total positivity test.

Can we use this idea to get k-positivity tests? Yes!

k-positivity Cluster Algebras

Total positivity seed where all variables = minors.



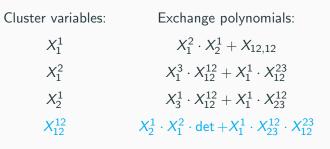
Frozen variables:

 $X_1^3 \qquad X_{12}^{23} \qquad X_3^1 \qquad X_{23}^{12} \qquad \det$

k-positivity Cluster Algebras

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Exchange polynomial uses minor of order $> k \implies$ freeze variable.



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k-positivity Cluster Algebras

Total positivity seed where all variables = minors.

Exchange polynomial uses minor of order $> k \implies$ freeze variable. Delete variables whose minors are "too big".

Cluster variables: Exchange polynomials:

X_1^1	$X_1^2 \cdot X_2^1 + X_{12,12}$
X_{1}^{2}	$X_1^3 \cdot X_{12}^{12} + X_1^1 \cdot X_{12}^{23}$
X_{2}^{1}	$X_3^1 \cdot X_{12}^{12} + X_1^1 \cdot X_{23}^{12}$
X_{12}^{12}	$X_2^1 \cdot X_1^2 \cdot \det + X_1^1 \cdot X_{23}^{12} \cdot X_{12}^{23}$

Frozen variables:

 $X_1^3 \qquad X_{12}^{23} \qquad X_3^1 \qquad X_{23}^{12} \qquad \det$

The *test cluster* of a seed is the extended cluster, but with more minors added until we have n^2 which combined give a *k*-positivity test. These extra *test variables* are the same for all seeds in the cluster algebra.

Example

Restricted *n*-initial minors seed + missing solid minors of order k = the *k*-initial minors test.

Don't (in general) know how to choose test variables to get a valid k-positivity test. Some seeds can't be extended to give tests (of size n^2) at all!

The *exchange graph* has vertices = clusters, and edges between clusters with exchange relations connecting them.

Example

For n = 2 total positivity cluster algebra:

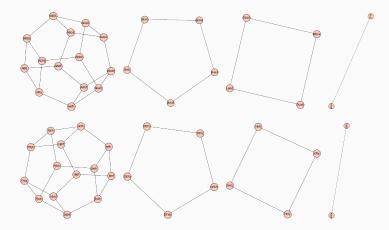


For 3×3 matrices, when we restrict exchanges to those only involving minors of size ≤ 2 , the exchange graph breaks into 8 components.

Only the two largest components provide actual 2-positivity tests.

These two components share 4 vertices that correspond to different total positivity tests but restrict to the same 2-positivity tests. We say that these 4 overlapping vertices form a "bridge" between the components.

Connected Components of 2-pos test graph for 3×3 matrix



Test Components

ABD ADh ABf Ade

Ide Jabd

Frozen variables: c,g,C,G,A Test variable: J

Frozen variables: c,g,C,G,J Test variable: A

k-essential minors

Definition

A minor is *k*-essential if there exists a matrix in which all other minors of size $\leq k$ are positive, while that minor is non-positive.

In other words, a k-essential minor is one which must be present in all k-positivity tests.

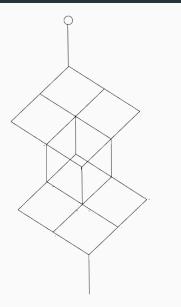
Conjecture

The k-essential minors are the corner minors of size < k, together with all solid k-minors.

So far, this conjecture has only been proven for the cases of $k \leq 3$.

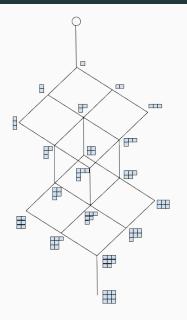
We also observe that in all known cases, a bridge involves switching the positions of an essential minor in the extended cluster with one outside it.

Connecting Tests



Although there are many choices to be made regarding the exact order in which some exchanges are made, we can generally speak of a natural family of paths linking the k-initial minors test to its antidiagonal flip. If we ignore non-bridge mutations and treat each connected component as a single vertex, we get a "bridge graph".

Connecting Tests



By the construction of the path, all involved bridges switch out a solid k-minor with a minor one entry down and to the left of it, yielding a total of $(n-k)^2$ distinct bridges, that we can represent as boxes in a $(n-k) \times (n-k)$ square. The components can thus be indexed by Young diagrams, with each box indicating a specific bridge that must be crossed to reach that component from the one including the k-initial minors.

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n = 5, k = 2

	2,2	2,3	1,3	1,4	1,5
	2,1	23,23	23, 34	12, 34	12,45
$ \longrightarrow$	3,1	23, 12	123,123	123, 234	123, 345
	4,1	34,12	234, 123	234, 234	234, 345
	5,1	45,12	345,123	345,234	345,345

			1,3	1,4	1,5
			23, 34		12,45
$ \longrightarrow \longrightarrow $				123, 234	
	5,2	45,23	234, 123	234, 234	234, 345
	5,1	45,12	345,123	345, 234	345,345

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